

\mathcal{N} -WL: A New Hierarchy of Expressivity for Graph Neural Networks

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k -Weisfeiler-Lehman (k -WL) hierarchy is a theoretical framework for graph isomorphism tests

./ but not practically useful when $k \geq 3$!

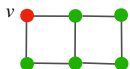
GIN = 1-WL [Xu et al., 2019]

Many expressive GNNs go beyond 1-WL

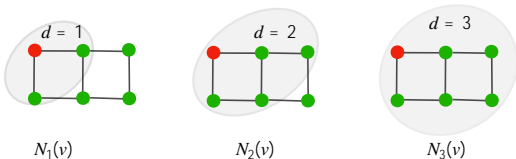
Question:

Is k -WL hierarchy a good yardstick for measuring expressivity of GNNs?

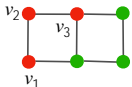
Colouring nodes



Local neighbourhood

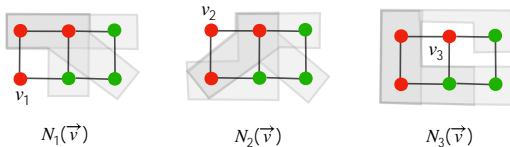


Colouring k -tuples

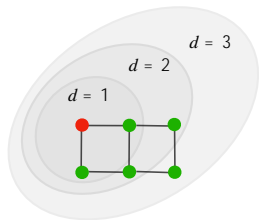


Global neighbourhood

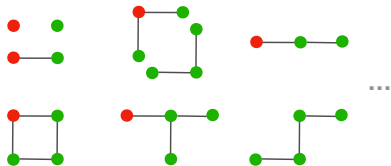
$$\vec{v} = \{v_1, v_2, v_3\}$$



\mathcal{N} -WL hierarchy computes node coloring via t -order induced subgraphs within d -hop neighbourhoods.



d -hop neighbourhoods



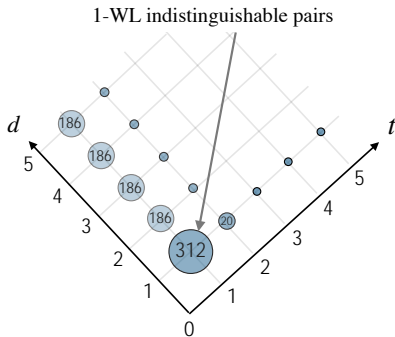
t -order induced subgraphs

A Simple Experiment

A graph isomorphism test on 312 pairs of simple graphs of 8 vertices:

None-or-all: none by 1-WL but all by 3-WL

Progressive: varying with d and t by \mathcal{N} -WL



Increasing the order of induced subgraphs, the expressive power increases
Not surprising

Theorem:
(Weak Hierarchy) $\mathcal{N}(t; d)\text{-WL} \subseteq \mathcal{N}(t+1; d)\text{-WL}$

Increasing the hops of neighbourhood, the expressive power may decrease
Surprising but can be proved

Theorem:
(Strong Hierarchy) $\mathcal{N}(t; d)\text{-WL} \subseteq \mathcal{N}(t+1; d)\text{-WL}$
 $\mathcal{N}(t; d)\text{-WL} \subseteq \mathcal{N}(t; d+1)\text{-WL}$

Induced connected subgraphs remain the same expressive power
Surprising but can be proved

Theorem:
(Equivalence) $\mathcal{N}^c(t; d)\text{-WL} = \mathcal{N}(t; d)\text{-WL}$

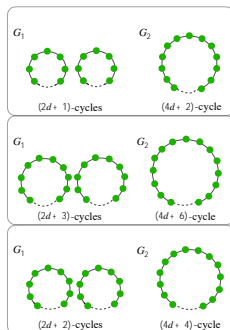
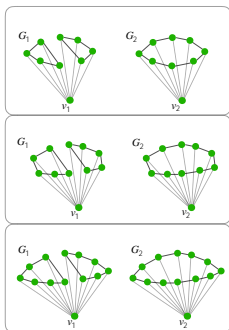
Main Ideas in Proofs (1)

Theorem:
(Strong Hierarchy)

$\mathcal{N}(t;d)$ -WL ($\mathcal{N}(t+1;d)$ -WL

$\mathcal{N}(t;d)$ -WL ($\mathcal{N}(t;d+1)$ -WL

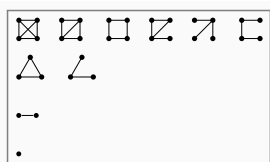
We prove strictness of hierarchies by constructing counterexample graphs.



Main Ideas in Proofs (2)

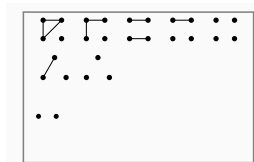
Theorem:
(Equivalence) $\mathcal{N}^c(t;d)$ -WL \iff $\mathcal{N}(t;d)$ -WL

Our proof is based on Kocay's Vertex Theorem [Kocay, 1982].



subgraph counts

implies
!



subgraph counts

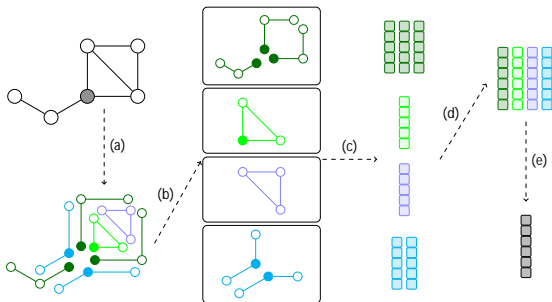
k-WL Hierarchy vs \mathcal{N} -WL Hierarchy

	k-WL	-k-LWL	(k;s)-LWL	(k;c)()-SETWL
#Coloured objects	n^k	n^k	subset(n^k ;s)	subset($\prod_{q=1}^k \binom{n}{q}$;c)
#Neighbour objects	$n \quad k$	$a \quad k$	$a \quad k$	$n \quad q$
Δ Coloured objects	k-tuples	k-tuples	k-tuples	k-sets
Δ Neighbour objects	k-tuples	k-tuples	k-tuples	...
Sparsity awareness	\times	\checkmark	\checkmark	

$\mathcal{N}(t;d)$ -WL	$\mathcal{N}^c(t;d)$ -WL
n	n
$\binom{a^d}{t}$	subset($\prod_{q=1}^t \binom{a^d}{q}$;1)
nodes	nodes
t-sets	t-sets
\times	\checkmark

Theorem: 1-WL $\mathcal{N}(1;1)$ -WL $\mathcal{N}^c(1;1)$ -WL

Graph Neighbourhood Neural Network (G3N) instantiates the ideas of \mathcal{N} -WL algorithms for graph learning.



$$h_u^{(l+1)} = \text{COMBINE } h_u^{(l)}; \text{AGG}_{(i;j)2l_t}^N \text{ } J_d \text{ } \text{AGG}_{S2S_u^{(l)}(i;j)}^T \text{ } \text{POOL}(S)$$



Kocay, W. L. (1982).

Some new methods in reconstruction theory.

In *Combinatorial Mathematics IX*, pages 89–114. Springer.



Xu, K., Hu, W., Leskovec, J., and Jegelka, S. (2019).

How powerful are graph neural networks?

In *International Conference on Learning Representations (ICLR)*.

Thank You