## Heuristic Search for Multi-Objective Probabilistic Planning

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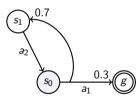
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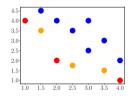
- stochastic shortest path problems (SSPs) are the defacto model for planning under uncertainty
  - compute a policy which maps states to actions
  - want optimal policy minimising the expected cost to reach the goal from an initial state



An SSP with a stochastic action  $a_1$ .

 heuristic search powerful for implicitly represented SSPs with large state spaces

- multi-objective stochastic shortest path problems (MOSSPs) generalise SSPs by exhibiting multiple objectives
  - aim to compute a <u>set</u> of non-dominated vector-valued policies with tradeoffs between objectives



orange+red=non-dominated vectors

- no existing heuristic search algorithms for MOSSPs
- algorithms that exist for MOMDPs do not carry over to MOSSPs: MOSSPs require additional assumptions

defining MOSSPs: a generalisation of MOMDPs

assumptions for convergence of MOVI for MOSSPs

heuristics for MOSSPs

heuristic search algorithms for solving MOSSPs: iMOLAO\*, MOLRTDP

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#### Definition (MOSSPs)

An MOSSP is a tuple  $P = (S, A, s_0, G)$ 

S is a set of states

A is a set of actions a which

maps certain states to a probability distribution of states, and

▶ has an associated cost vector  $\vec{C}(a) = [c_1, \dots, c_n] \in \mathbb{R}_{\geq 0}^n$ 

▶ 
$$s_0 \in S$$
 is an initial state,  $G \subseteq S$  is a set of goal states

Informally,

A solution is a compact representation of a set of non-dominated vector-valued policies.

 $\vec{u}$  dominates  $\vec{v}$  denoted  $\vec{u} \leq \vec{v}$  iff  $\vec{u}[i] \leq \vec{v}[i], i = 1, \dots, n$ 

\*Full definition: too many technical details and sub-definitions to fit into the talk

## MO Bellman Backup

- fundamental equations for solving (MO)SSPs and (MO)MDPs
- ▶ states assigned **Q** and **V** values; each a finite set of vectors  $\in \mathcal{P}_{<\omega}(\mathbb{R}^n_{\geq 0})$
- ▶  $\mathbf{Q}: S \times A \rightarrow \mathcal{P}_{<\omega}(\mathbb{R}^n_{\geq 0}) \simeq \mathsf{MO}$  expected cost to goal when executing *a* at *s*
- ▶  $V: S \to \mathcal{P}_{<\omega}(\mathbb{R}^n_{>0}) \simeq MO$  expected cost to goal from  $s \to minimum$  of Q values

For 
$$s \in S \setminus G$$
:  

$$\mathbf{V}^{t+1}(s) = \operatorname{CCS}\left(\bigcup_{a \in A} \mathbf{Q}^{t+1}(s, a)\right)$$

$$\mathbf{V}^{t+1}(g) = \{\vec{0}\}$$

$$\mathbf{V}^{t+1}(s, a) = \{\vec{C}(a)\} \oplus \left(\bigoplus_{s' \in S} P(s'|s, a)\mathbf{V}^{t}(s')\right)$$
red: CCS;  
orange+red: PCS

- $\blacktriangleright$  CCS is the convex hull of a Pareto coverage set  $\simeq$  MO generalisation of min
- ▶  $\mathbf{V} \oplus \mathbf{U} = \{\vec{u} + \vec{v} \mid \vec{u} \in \mathbf{U}, \vec{v} \in \mathbf{V}\}$  is setwise sum of vectors
- $\bigoplus$  generalises  $\oplus$  to several sets
- $\lim_{t \to \infty} \mathbf{V}^t = \mathbf{V}^*$  [White, 1982; Barrett and Narayanan, 2008]

# MO Value Iteration

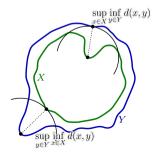
#### Algorithm: MOVI

**Data:** MOSSP problem  $P = (S, A, s_0, G)$ , initial values  $\mathbf{V}(s)$  for each state s (default to  $\mathbf{V}(s) = \{\vec{0}\}$ ), and consistency threshold  $\varepsilon$ . 1 while  $\max_{s \in S} res(s) < \varepsilon$  do for  $s \in S$  do 2 if  $s \in G$  then  $V_{\text{new}}(s) \leftarrow \{\vec{0}\}$ : 3 else  $V_{new}(s) \leftarrow$ 4 MOBellmanBackup(s);  $\mathsf{res}(s) \leftarrow D(\mathbf{V}, \mathbf{V}_{\mathsf{new}})$ 5  $V \leftarrow V_{new}$ 7 return V

same as (single-objective) Value Iteration except

- 1. MO backups instead of SO backups
- 2. MO residual defined using Hausdorff metric (for finite sets of points):

$$D(\mathbf{U}, \mathbf{V}) = \max \left\{ \max_{\substack{\vec{u} \in \mathbf{U} \\ \vec{v} \in \mathbf{V}}} \min_{\vec{v} \in \mathbf{U}} d(\vec{u}, \vec{v}), \\ \max_{\substack{\vec{u} \in \mathbf{V} \\ \vec{v} \in \mathbf{U}}} \min_{\vec{v} \in \mathbf{U}} d(\vec{u}, \vec{v}) \right\}$$



#### defining MOSSPs: a generalisation of MOMDPs

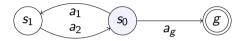
#### assumptions for convergence of MOVI for MOSSPs

heuristics for MOSSPs

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Assumptions for convergence of MOVI for MOSSPs

Consider the  $\ensuremath{\mathsf{MOSSP}}$ 



MOSSP with action costs  $\vec{C}(a_1) = [1,0], \vec{C}(a_2) = [1,0], \vec{C}(a_g) = [0,1].$ 

• 
$$\pi_1: S \to A, s_0 \mapsto a_g$$
 is a proper policy

•  $\pi_2: S \to A, s_0 \mapsto a_1, s_1 \mapsto a_2$  is an improper policy

- ▶ MOVI does not detect this;  $V(s_0) = \{[\infty, 0], [0, 1]\}$  in the limit
- $\blacktriangleright$  VI able to detect and prune improper SSP policies which have  $\infty$  cost

# Assumptions for convergence of MOVI for MOSSPs

Strong improper policy assumption

- reachability assumption
  - ▶  $\forall s \in S$ ,  $\exists$  proper policy at s
- $\blacktriangleright\,$  all improper policies cost  $\vec{\infty}$

- MOVI + strong assumption is sound and complete
- implies cycle costs have nonzero components (not easy to guarantee)
- one method: all action costs have nonzero components
  - cannot model independent costs
  - e.g. navigation domain: load, unload consumes 0 fuel, 1 time

#### Weak improper policy assumption

- reachability assumption
- exists a bound  $\vec{b}$  s.t.  $\forall s \in S$ 
  - for all proper  $\pi$ ,  $V^{\pi}(s) \preceq \vec{b}$
  - for all improper  $\pi$ ,  $V^{\pi}(s) \not\preceq \vec{b}$

• 
$$\vec{u} \leq \vec{v}$$
 iff  $\vec{u}[i] \leq \vec{v}[i], i = 1, \dots, n$ 

- modified MOVI + weak assumption is sound and complete
- $\blacktriangleright$  can derive or estimate upper bound  $\vec{b}$

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- heuristic search powerful for (optimal) planning
- does not require enumerating the whole potentially exponential search space
- heuristic search algorithms:

	SO	MO
determinstic	A*, BiA* etc.	NAMOA*
stochastic	(i)LAO*, LRTDP etc.	our contributions

## **MOSSP** Heuristic

- admissible heuristics in single-objective (SO) deterministic case guarantee optimality with A\*
- ▶ near optimality for SSPs with (i)LAO\* and LRTDP in finitely many iterations
  - optimality with possibly infinitely many iterations
- ▶ in SO case, heuristic for a state is a scalar  $h(s) \in \mathbb{R}_{\geq 0}$
- ▶ in MO case, heuristic for a state is a <u>set</u> of vectors  $\mathbf{H}(s) \subset \mathbb{R}_{\geq 0}^n$

#### Definition (admissible MO heuristic)

**H** is admissible if  $\forall s \in S \setminus G$ , for all  $\vec{v} \in \mathbf{V}^*(s)$  there exists  $\vec{u} \in \mathbf{H}(s)$  such that  $\vec{u} \leq \vec{v}$  where  $\mathbf{V}^*$  is the optimal value function, and  $\forall g \in G, \mathbf{H}(g) = \{\vec{0}\}$ .

▶ same definition as in deterministic MO case [Mandow and Pérez-de-la-Cruz, 2010]

## Domain independent MOSSP Heuristics

- ▶ no heuristic: the zero heuristic defined by  $\mathbf{H}_{zero}(s) = {\vec{0}}$
- can use SO heuristics
  - *ideal point heuristic*; apply an SO heuristic  $h_i$  to each objective in isolation:

$$\mathbf{H}_{ideal}(s) = \{[h_1(s), \dots h_n(s)]\}$$

- e.g.  $\mathbf{H}_{ideal}^{max}$ ,  $\mathbf{H}_{ideal}^{pdb2}$ ,  $\mathbf{H}_{ideal}^{pdb3}$  [Bonet and Geffner, 2001; Klößner and Hoffmann, 2021]
- can use deterministic MO heuristics
  - construct determinised problem by replacing each probabilistic effect with a deterministic action
  - e.g. H<sup>comax</sup><sub>mo</sub>, H<sup>pdb2</sup><sub>mo</sub>, H<sup>pdb3</sup><sub>mo</sub> [Geißer et al., 2022]
- new abstraction heuristics
  - combine values from solving smaller projections of the problem
  - considers both MO and stoch. features of MOSSPs
  - e.g.  $\mathbf{H}_{mossp}^{pdb2}$ ,  $\mathbf{H}_{mossp}^{pdb3}$
- other existing SO and MO heuristics can be extended in above ways

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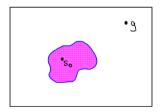
## iMOLAO\*

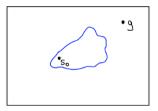
gradually build partial solution via heuristic search

#### Algorithm: iMOLAO\*

Data: MOSSP problem  $P = (S, A, s_0, G)$ , heuristic **H**. and consistency threshold  $\varepsilon$ 1  $\mathbf{V} \leftarrow \mathbf{H}; \ \Pi \leftarrow \emptyset; \ F \leftarrow \{s_0\}; \ I \leftarrow$  $\emptyset: N \leftarrow \{s_0\}$ 2 while  $((F \cap N) \setminus G \neq \emptyset) \land (\max_{s \in N} res(s) < \varepsilon)$ do  $F = \emptyset$ 3  $N \leftarrow \text{postorderTraversalDFS}(s_0, \Pi)$ for  $s \in N$  in the computed order do 5  $V(s) \leftarrow MOBellmanBackup(s)$ 6  $\Pi(s) = \text{getActions}(s, \mathbf{V})$ 7 if  $s \notin I$  then  $F = F \cup \{s\}$ ; 8  $I = I \cup \{s\}$ q 10 return V

1. collect states N of the current best partial solution  $\Pi$  in postorder traversal DFS order





3. update new current partial solution  $\Pi;$  then return to step 1 and repeat until covergence criteria is met

2. run MO Bellman Backups once on each state in the collected order

# MOLRTDP

#### Algorithm: MOLRTDP

```
Data: MOSSP problem P = (S, A, s_0, G), heuristic
            H, and consistency threshold \varepsilon
    procedure MOLRTDP(P, \varepsilon, \mathbf{H})
        V \leftarrow H
        while \neg s_0. solved do
 2
             visited \leftarrow \emptyset
             s \leftarrow s_0
             while ¬s solved do
                  visited.push(s)
 6
                  if s \in G then break :
                  V(s) \leftarrow MOBellmanBackup(s)
 8
                  a \leftarrow \text{sampleUnsolvedGreedyAction}(s)
 q
                  s \leftarrow \text{sampleUnsolvedNextState}(s, a)
10
             while ¬visited.empty() do
11
                  s \leftarrow visited.pop()
12
                  if ¬checkSolved(s) then break ;
13
        return V
14
```

```
routine checkSolved(s)
     rv \leftarrow true; open \leftarrow \emptyset; closed \leftarrow \emptyset
    if \neg s. solved then open. push(s) :
    while \neg open.empty() do
         s \leftarrow open.pop()
         if res(s) > \varepsilon then
               rv \leftarrow false
              continue
         for a \in getActions(s, V) do
              for s' \in \text{successors}(s, a) do
                   if \neg s'.solved \land s' \notin open \cup closed
                     then open.push(s'):
    if rv then for s \in closed do s.solved = true:
    else
         while closed \neq \emptyset do
              s \leftarrow closed.pop()
               V(s) \leftarrow MOBellmanBackup(s)
    return rv
```

2

3

Δ

5

6

8

q

10

11

12

13

14

15

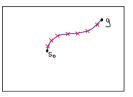
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# MOLRTDP

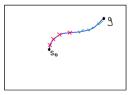
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```
Data: MOSSP problem P = (S, A, s_0, G),
            heuristic H, consistency threshold \varepsilon
    procedure MOLRTDP(P, \varepsilon, \mathbf{H})
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         visited \leftarrow \emptyset
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         while ¬s_solved do
            visited.push(s)
            if s \in G then break :
            V(s) \leftarrow MOBellmanBackup(s)
            a \leftarrow \text{sampleUnsolvedGreedvAction}(s)
 a
            s \leftarrow \text{sampleUnsolvedNextState}(s, a)
10
         while \neg visited.emptv() do
11
            s \leftarrow visited.pop()
12
            if ¬checkSolved(s) then break ;
13
      return V
14
```

 a state is 'solved' if its own and all its ancestors' values under the current partial solution have converged



1. random greedy walk to the goal on unsolved states if possible, while performing backups along the way



2. update whether states are solved in reverse order of the trialled walk using checkSolved; return to step 1 and repeat until  $s_0$  is solved

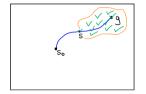
# MOLRTDP - checkSolved

```
routine checkSolved(s)
      rv \leftarrow true; open \leftarrow \emptyset; closed \leftarrow \emptyset
      if \neg s.solved then open.push(s);
 2
      while \neg open.empty() do
         s \leftarrow open.pop()
         if res(s) > \varepsilon then
            rv ← false
 6
            continue
         for a \in getActions(s, V) do
 8
            for s' \in \text{successors}(s, a) do
 a
               if \neg s'.solved \land s' \notin open \cup closed
10
                 then open.push(s');
      if ry then for s \in closed do
11
        s.solved = true:
      else
12
         while closed \neq \emptyset do
13
            s \leftarrow closed.pop()
14
            V(s) \leftarrow MOBellmanBackup(s)
15
16
      return rv
```

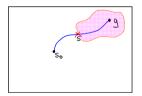
 a state is 'solved' if its own and all its ancestors' values under the current partial solution have converged

2.1. collect all greedy ancestor states of

s under the current partial solution



 $2.2.a.\ \mbox{if}$  values for all ancestor states have converged, set the state and its ancestors to be solved



2.2.b. else run backups on the state and all its ancestors

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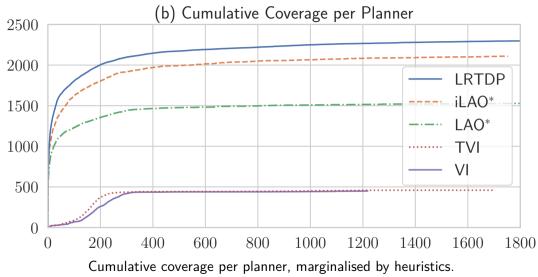
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## Experimental setup

- ▶ 5 solvers (MO prefix omitted): VI, TVI, LAO\*, iLAO\*, LRTDP
- ▶ 9 heuristics: zero, H<sup>max</sup><sub>ideal</sub>, H<sup>comax</sup>, H<sup>pdb2</sup><sub>ideal</sub>, H<sup>pdb3</sup><sub>ideal</sub>, H<sup>pdb2</sup><sub>mo</sub>, H<sup>pdb3</sup><sub>mossp</sub>, H<sup>pdb3</sup><sub>mossp</sub>, H<sup>pdb3</sup><sub>mossp</sub>
- ▶ 30 minute timeout, single CPU core, 4GB memory
- ▶ 7 domains, mix of MO and probabilistic interesting domains
- 610 total problems
- ▶  $5 \times 9 \times 610 = 27450$  possible experimental configurations

## Best planner



## Best heuristic

Heuristic	Normalised coverage	-	Heuristic	Unnormalised coverage
$\mathbf{H}_{mossp}^{pdb3}$	19.1	-	$H^{pdb3}_{mossp}$	893.5
$\mathbf{H}_{mossp}^{pdb2}$	17.9		$\mathbf{H}_{mossp}^{pdb2}$	893.3
$\mathbf{H}_{mo}^{pdb3}$	17.4		$\mathbf{H}_{mo}^{pdb3}$	871.2
$\mathbf{H}_{mo}^{pdb2}$	17.1		$\mathbf{H}_{ma}^{pdb2}$	851.8
$\mathbf{H}_{ideal}^{max}$	14.9		$\mathbf{H}_{ideal}^{pdb3}$	768.7
$\mathbf{H}_{ideal}^{pdb3}$	14.7		H <sup>max</sup> ideal H <sup>pdb2</sup>	755.2
$\mathbf{H}_{ideal}^{pdb2}$	14.4		$\mathbf{H}_{ideal}^{pdb2}$	737.2
blind	12.1		blind	572.9
$\mathbf{H}_{mo}^{comax}$	10.7	_	$\mathbf{H}_{mo}^{comax}$	504.5

marginalise by planner

- normalisation is done over domain due to uneven number of problems
- ranking of top 3 and bottom 3 heuristics same

## Heuristic accuracy

	Critical Path			Abstractions					
	blind	H max ideal	<b>Н</b> сотах то	H <sup>pdb2</sup> ideal	Н <sup>рdb3</sup> ideal	$H_{mo}^{pdb2}$	H <sup>pdb3</sup> mo	H <sup>pdb2</sup> mossp	H pdb3 mossp
SAR-4	100	97	44	93	92	44	44	38	26
SAR-5	100	97	45	93	92	45	45	39	28
ExBw-2d	100	59	24	52	45	52	45	52	45
ExBw-3d	100	58	22	52	44	52	44	52	44
Tireworld	100	100	68	100	100	68	68	57	12
VisitAll	100	100	54	100	100	61	53	22	12
VisitAllTire	100	100	50	100	100	66	45	66	45

- mean relative error (%) of heuristic value at the initial state relative to the optimal value for solved instances
- error calculated by directed Hausdorff metric difference divided by the norm of the largest vector of the optimal value:

 $\max_{\vec{v} \in \mathbf{V}} \min_{\vec{u} \in \mathbf{H}} d(\vec{v}, \vec{u}) / \max_{\vec{v} \in \mathbf{V}} \|\vec{v}\|$ 

 $\implies$  MO features more important than stoch. features to capture in a heuristic

Thank you for your attention!